



# Compression of radio astronomy data flows

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# Motivation

- Data transport and storage are costly, particularly at the scale of the SKA.
- Compression reduces data volumes that must be transported/stored, potential cost reductions follow.

(It comes with its own costs/overheads of course!)

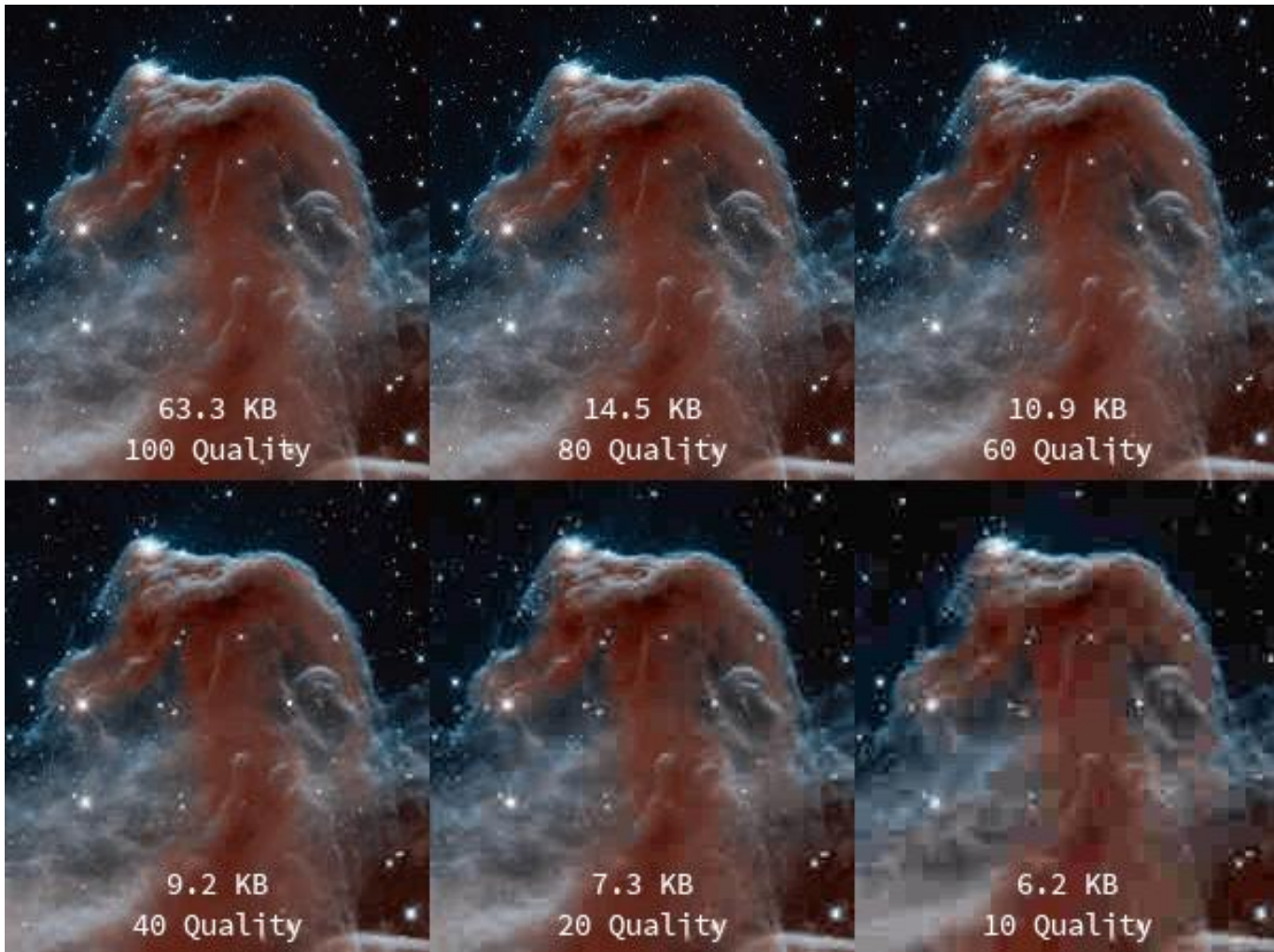
# The general scheme



# 2 types of compression

- Lossless – can completely recover the original information from the compressed data
- Lossy – information lost during the compression process is irretrievable

# Lossy JPEG compression



# Lossless, simple example

## Run Length Logic (RLL)

- Original data string = 55555555555555555555
- RLL string = 555#17
- Send 6 characters instead of 20!

# Data source characteristics

Need to know;

- Source Alphabet = set of all data symbols/states the source can generate

$$A = \{a, b, c, d, \dots, z\}$$

$$A = \{00, 01, 10, 11\}$$

- Symbol probability distribution

$$P = \{p(a), p(b), \dots, p(z)\}$$

$$P = \{p(00), p(01), p(10), p(11)\}$$

# Some information theory

- $I(a_i) = \log \left( \frac{1}{P(a_i)} \right) = -\log(P(a_i))$

- If base of log = 2 then unit = the **bit**

$$I(a_i) = -\log_2(P(a_i)) \text{ bits}$$

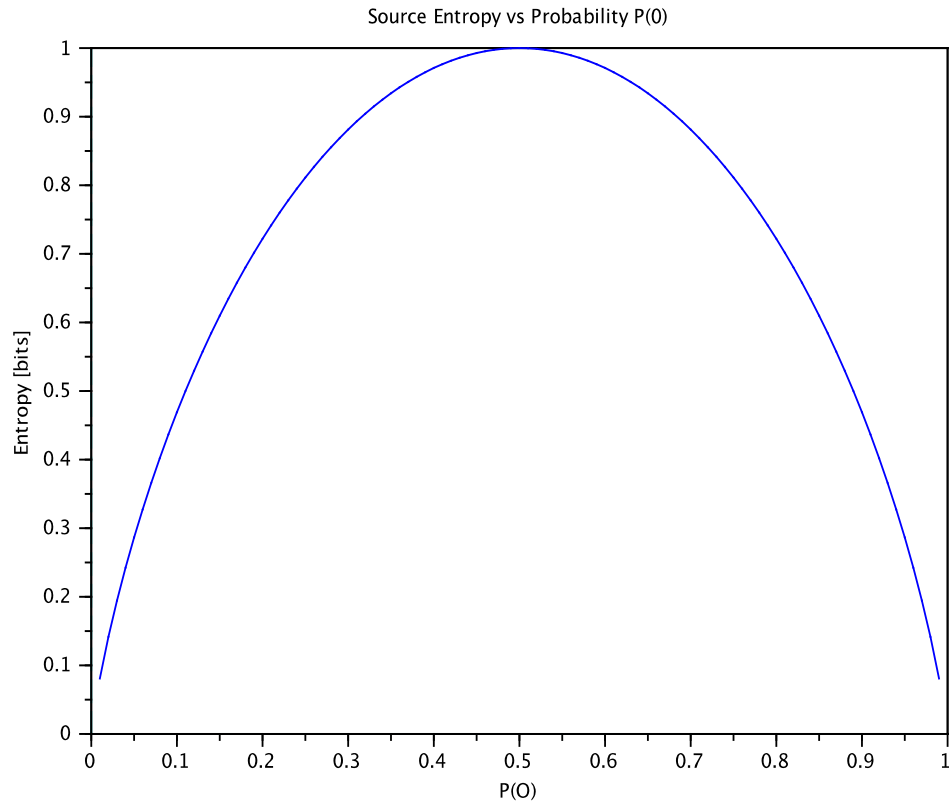
- Source Entropy = average no of bits per symbol

$$H(A, P) = \sum_{i=1}^N P(a_i) I(a_i) = -\sum_{i=1}^N P(a_i) \log_2(P(a_i))$$

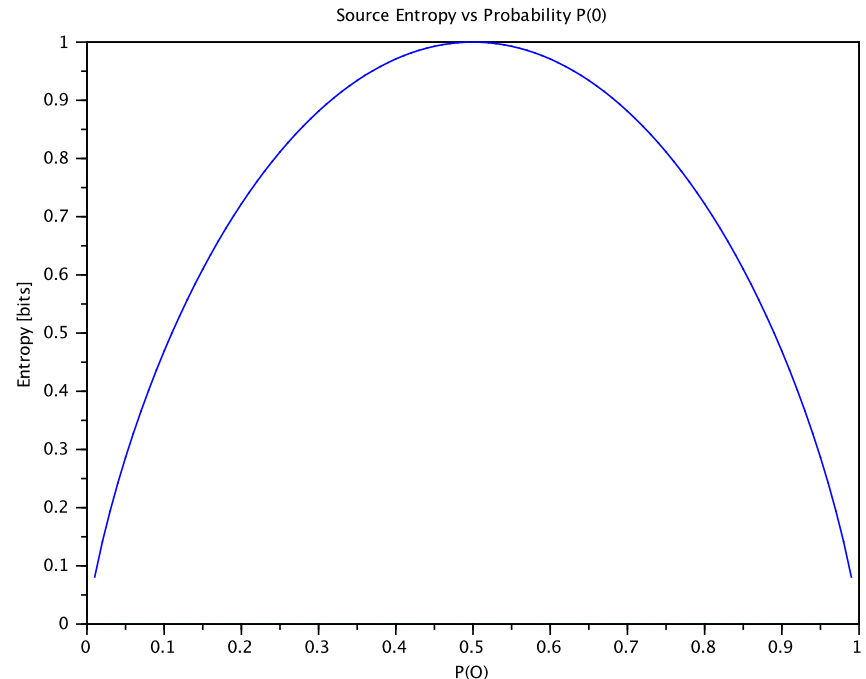


# Entropy in 1 bit case

- $A = \{0, 1\}$ ,  $P = \{p(0), p(1)\}$ ,  $p(0)+p(1)=1$



- Entropy is maximised when symbol emission is equiprobable!
- Entropy reduces as the distribution moves from the equiprobable state!

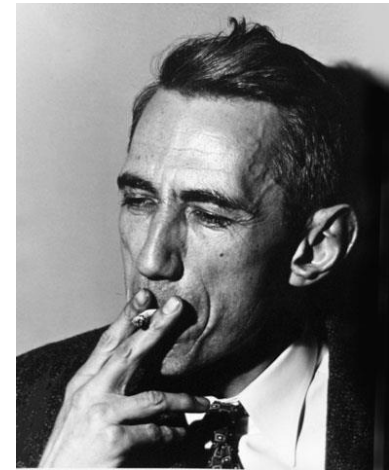


# Shannon lossless source coding theorem

$$H(A, P) \leq L \leq H(A, P) + \frac{1}{N}$$

- $L$  = minimum possible length of (losslessly) compressed data string
- For  $N \rightarrow \infty$  then

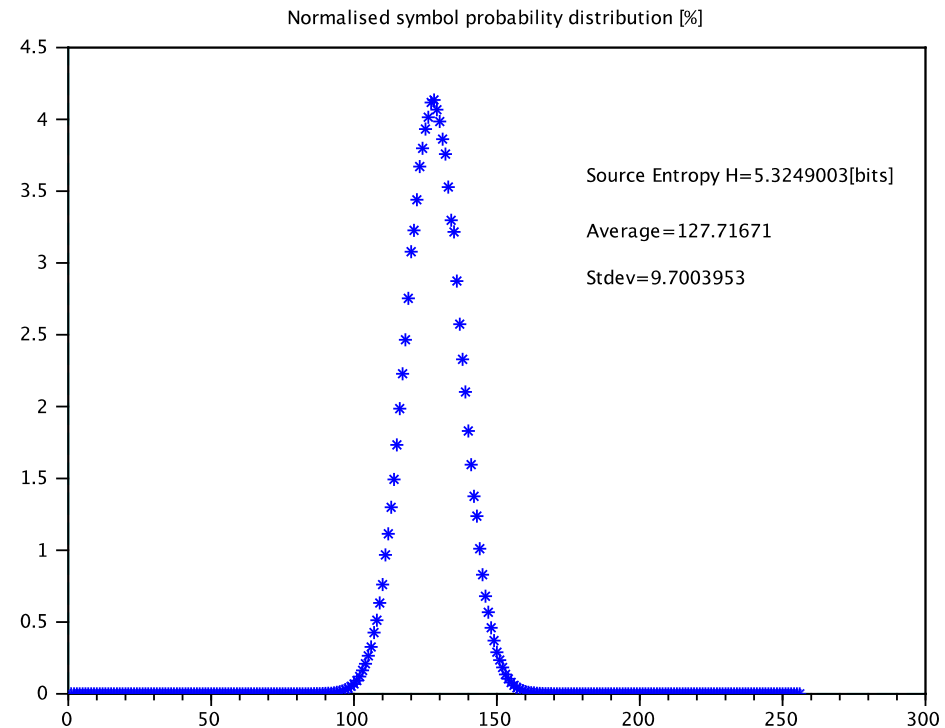
$$L = H(A, P)$$



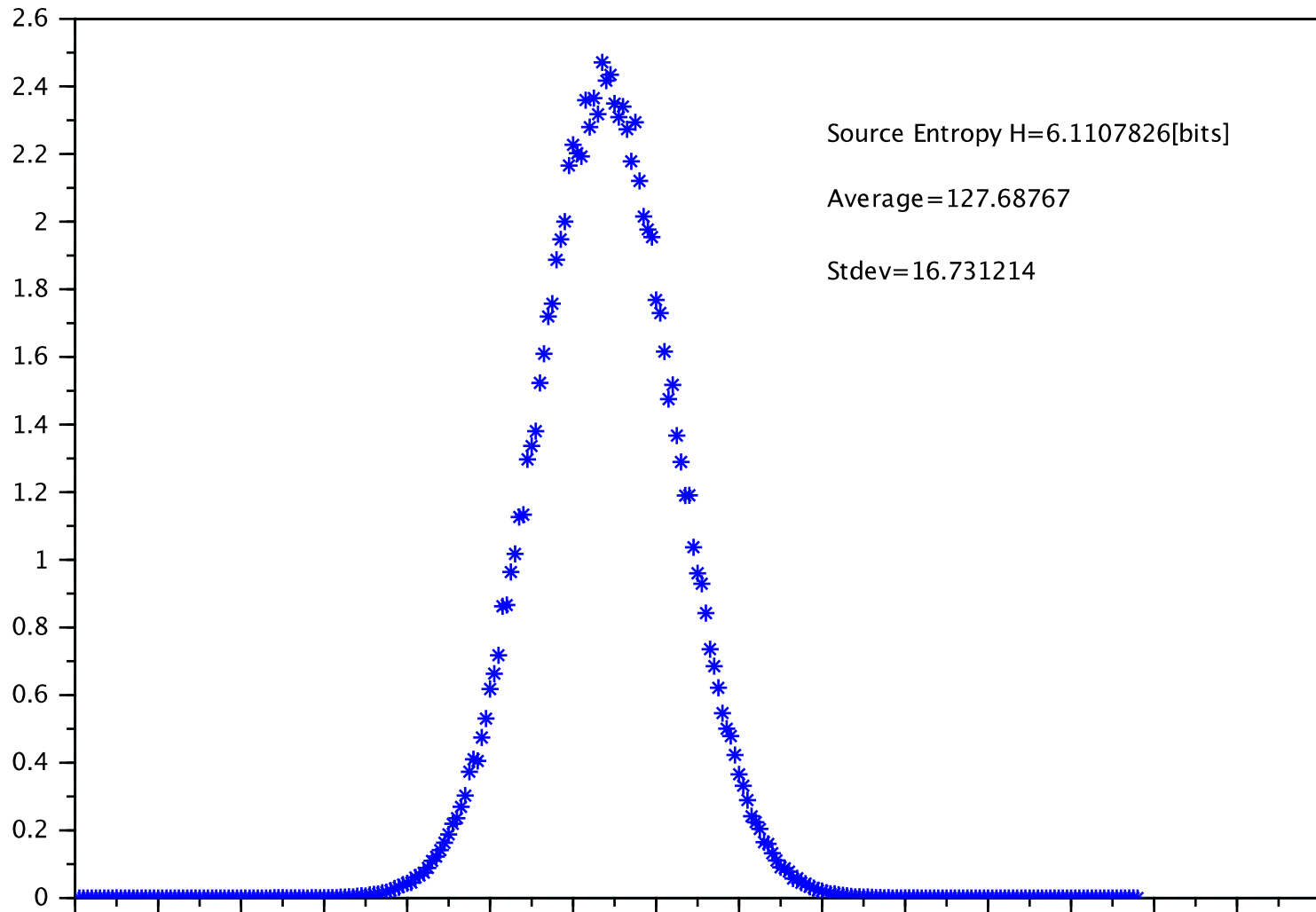
Died 2001 aged 85 years

# Radio Astronomy data sources

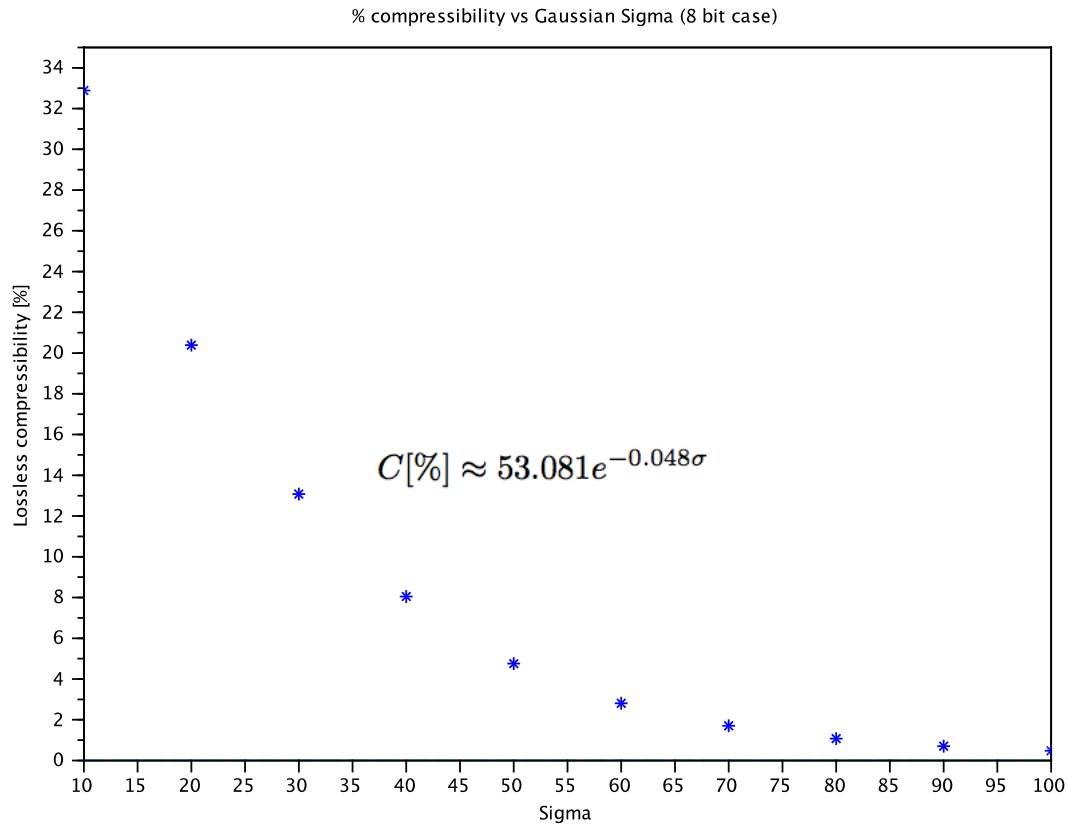
- Typically have a gaussian probability distribution- **non equiprobable distribution of states!**
- 8 bit example (real)



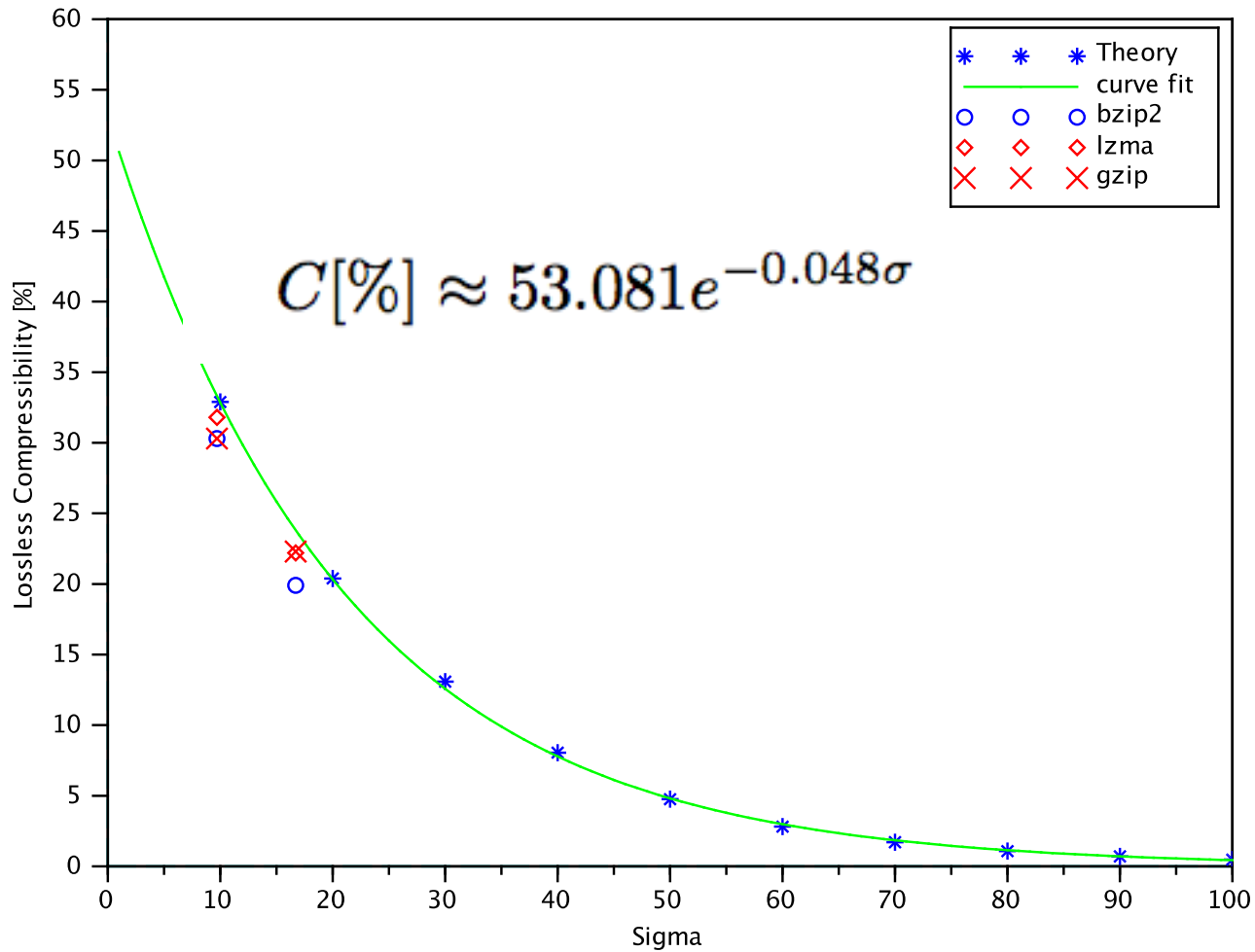
# Another file



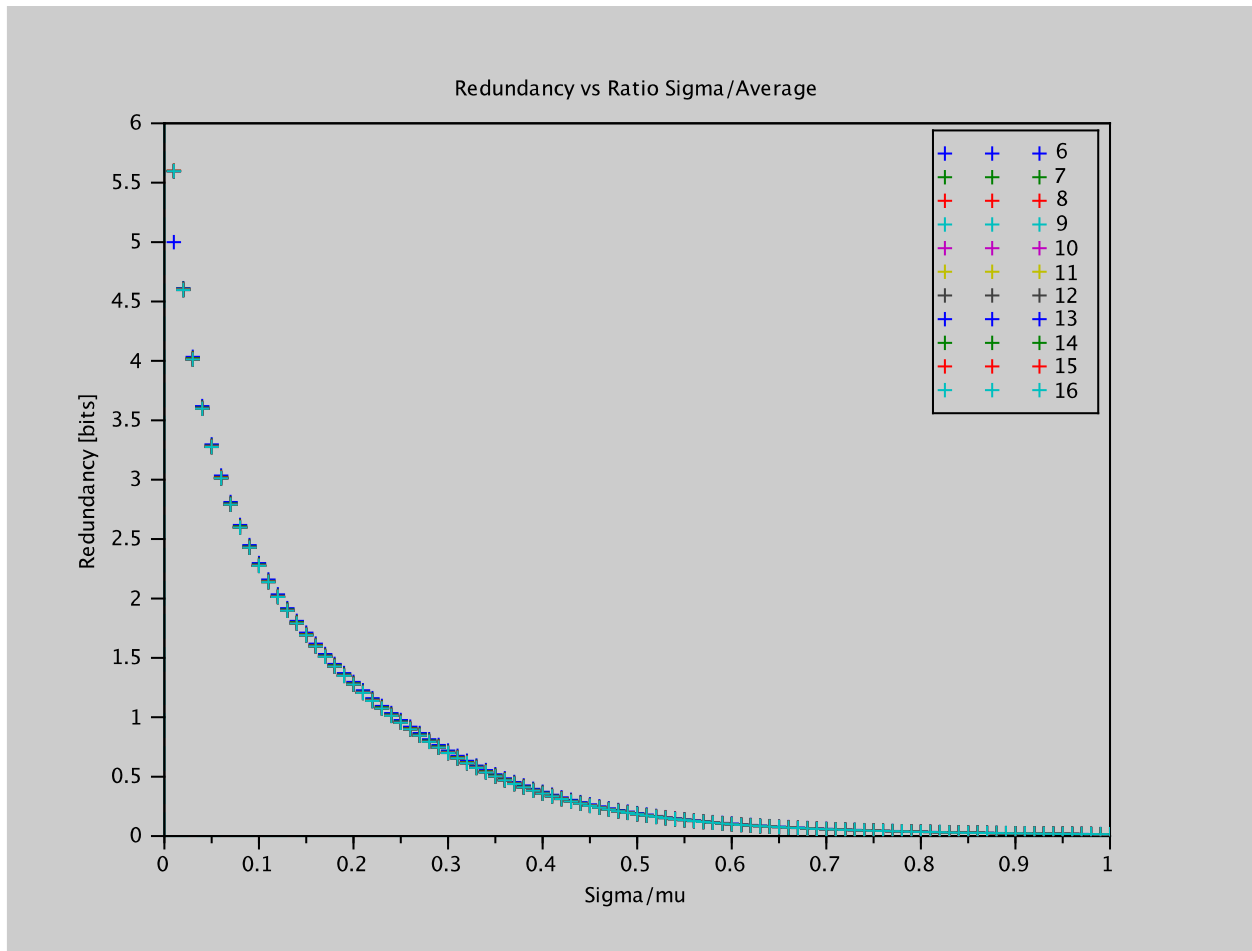
# Results from simulations of 8 bit gaussian distributed data



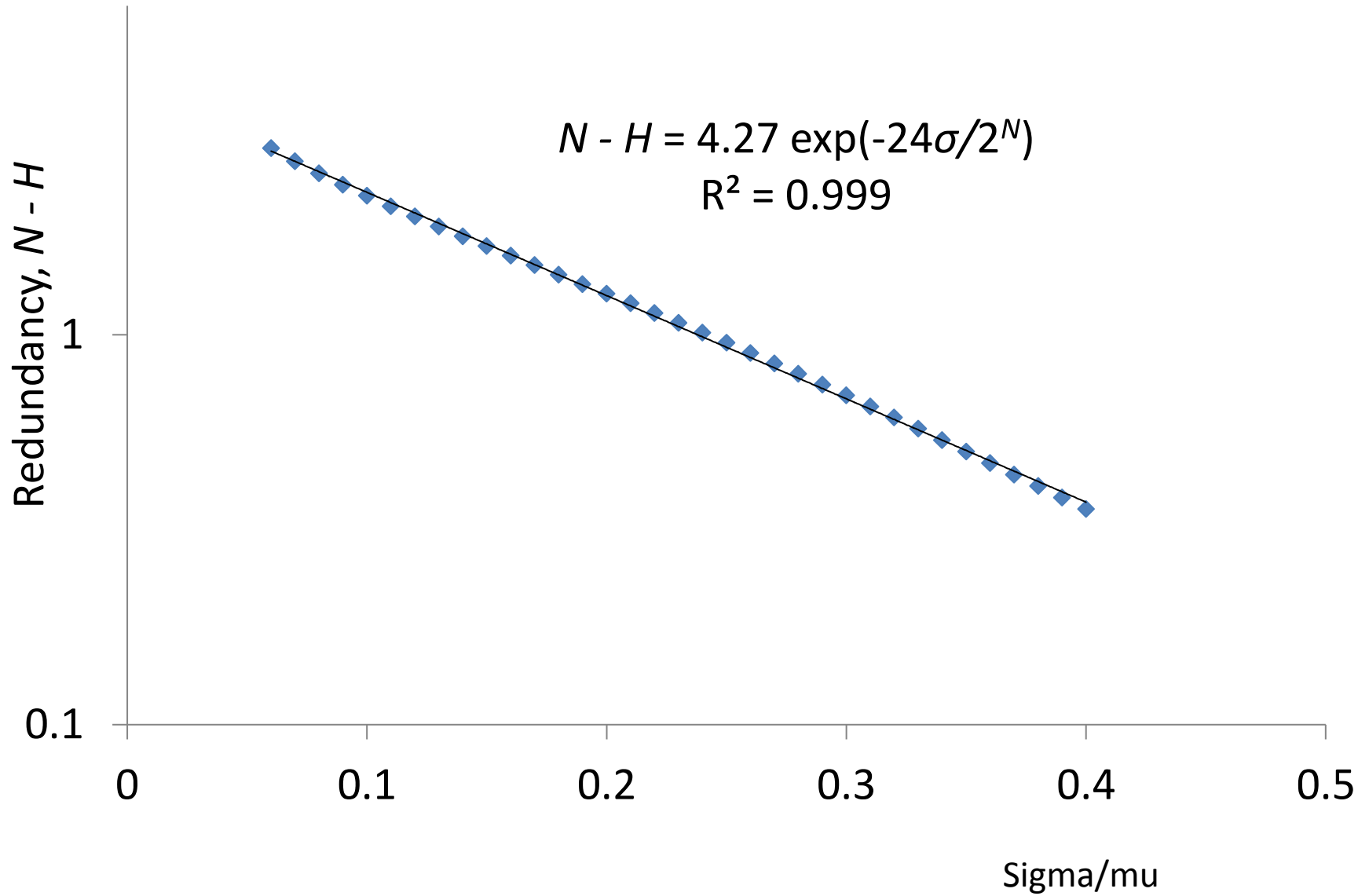
# Theory vs practice (8 bit case)



# Generalisation to N bit digitisation







# Proviso's / assumptions

- Gaussian distribution
- Memoryless source; probability of observing symbol  $a_i$  is independent of any previously emitted symbols
- Non memoryless state requires a different definition of entropy and offers yet more scope for compression

$$H(A, P) = \sum_{i=1}^N P(a_i) \sum_{j=1}^N P(a_j|a_i) \log_2(P(a_j|a_i))$$